

ISMP 2012,
Berlin

Alan Holland
and Barry
O'Sullivan

Introduction

Mechanism
Design

Potential
E-Policy
Demonstrator
Application

Research
Challenges:
Allocation

Optimising the economic efficiency of monetary incentives for renewable energy investment

Alan Holland and Barry O'Sullivan

Cork Constraint Computation Centre
University College Cork
Ireland

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Problem Description

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- Government policy: increase renewable energy production.
- Limited budgets.
- Lack of information regarding willingness of population to share costs.
- Instrument to support policy: grant aid.
- **Primary challenge: who should get what level of grant aid.**
- Secondary challenge: how to split a budget among competing technologies.

Game Theory

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- A formal way to analyze interaction among rational agents who behave strategically.
- Economic modelling tool.
- People behave in a **selfish** manner that maximizes their own utility.

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- A formal way to analyze interaction among rational agents who behave strategically.
- Economic modelling tool.
- People behave in a **selfish** manner that maximizes their own utility.
- Solution concept: Nash equilibrium
- Outcome is stable (no agent has an incentive to *unilaterally* deviate)

Inverse Game Theory

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- game of *private* information
- center chooses the payoff structure
- agent's "type" $\theta \in \Theta$
- outcome o consists of an allocation and payoff
 $o(\theta) = \{x(\theta), p(\theta)\}, x \in X, p \in P$

Designing An Economic Game

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- Inverse Game Theory
- Design rules
- selfish actions lead to socially desirable outcome

Designing An Economic Game

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Allocation Scheme

Who gets what.

Payment Scheme

What they pay to the Center in
exchange.

Solar Grant Problem

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- Altruistic central planner.
- Finite budget b to subsidize renewable energy micro generation.
- Set of self-interested agents I (seeking subsidies).
- Private information held by agent i^1
- Common knowledge: price to acquire and install r
 - Pitch of roof p_i
 - Orientation o_i
 - Latitude l_i
 - **Value of cashflow stream v_i**

¹Smart phone applications available for this purpose (e.g. Pitch Gauge) 

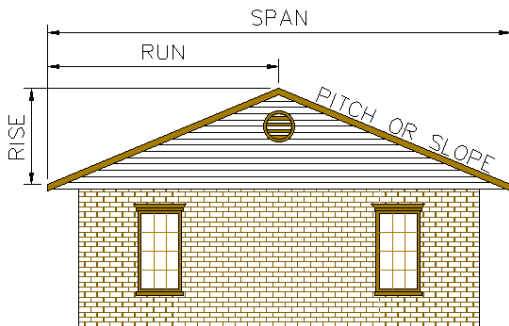


Figure: Roof pitch.

$\langle p, o, l, v \rangle$

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Figure: Orientation.

$\langle p, o, l, v \rangle$

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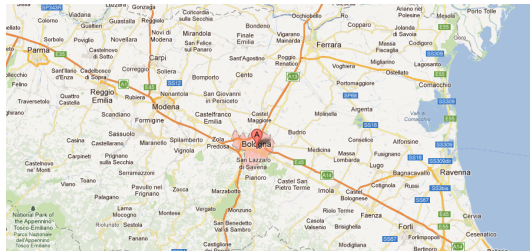


Figure: Latitude

Smart Phone: $\langle p, o, l, v \rangle$

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Figure: Data Capture

Cost

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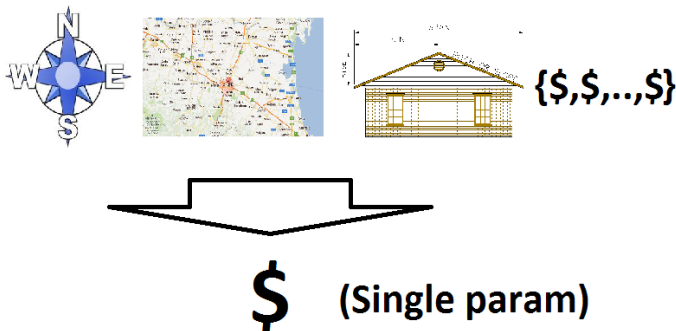


Figure: Maximum Cost an agent is willing to pay

Goal of Center

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- **Minimize maximum cost for any agent**
- Incentivise truthful revelation of private data
- Allocate grants to those best placed to generate solar energy cost effectively
- Respect budgetary constraints.

Problem and Rationale

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Challenges:
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- Makespan minimization problem ($Q||C_{max}$)
- NP-complete
- allocate device acquisition and hosting responsibility (jobs) across houses (machines)
- private cost associated with acceptance of that job.
- inconvenience is bounded as tightly as possible.

Greedy Allocation Algorithm

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- 1 order the devices from most expensive to cheapest
- 2 Greedily assign each to the household that minimizes the max cost imposition.
- 3 A 2-approximation that is **non-monotone**.
- 4 Leads to strategic manipulability

Non-monotone Allocation Failure Example

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- 3 devices $\{d_1, d_2, d_3\}$ and 2 house-owners $\{h_1, h_2\}$
- power of each device: $\phi_1=10W$, $\phi_2 = \phi_3 = 9 + \epsilon W$.
- price for each device: \$60 (common to all).
- Each house-owner has private value of \$5/W and $(5 - \epsilon)\$/W$, resp.
- greedy algorithm $d_1 \rightarrow h_1$, $d_2 \rightarrow h_2$ and $d_3 \rightarrow h_3$
- Cost $c_1 = 60 - (10 \times 5) = \50 ,
 $c_2 = 2 \times (60 - (45 - 4\epsilon - \epsilon^2)) \approx \30 .
- Note: if v_2 increases to $5 + \epsilon$, he loses the second device.

Solution

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- Complete search to ensure optimality
- Intractable: potentially large problem instances
- Compromise: approximation scheme with guarantees of monotonicity.

Approach: Monotone Algorithm

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- 1 Random 3-approximation [Kovacs 2005]
- 2 Randomized rounding of partial allocations
- 3 Implementable within a truthful mechanism (critical payment scheme).

Payment Scheme

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- Critical payments
- The minimum cost declaration to be awarded that item.
- Implementable within a truthful mechanism.

Future Work

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Challenges:
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- Empirical study of expected outcomes using simulations
- Pilot Subsidy Auction
- Emilia Romagna region of Italy